

# Virtual Consensus in Delos

Balakrishnan et al., OSDI '2020

# Consensus-based database replication

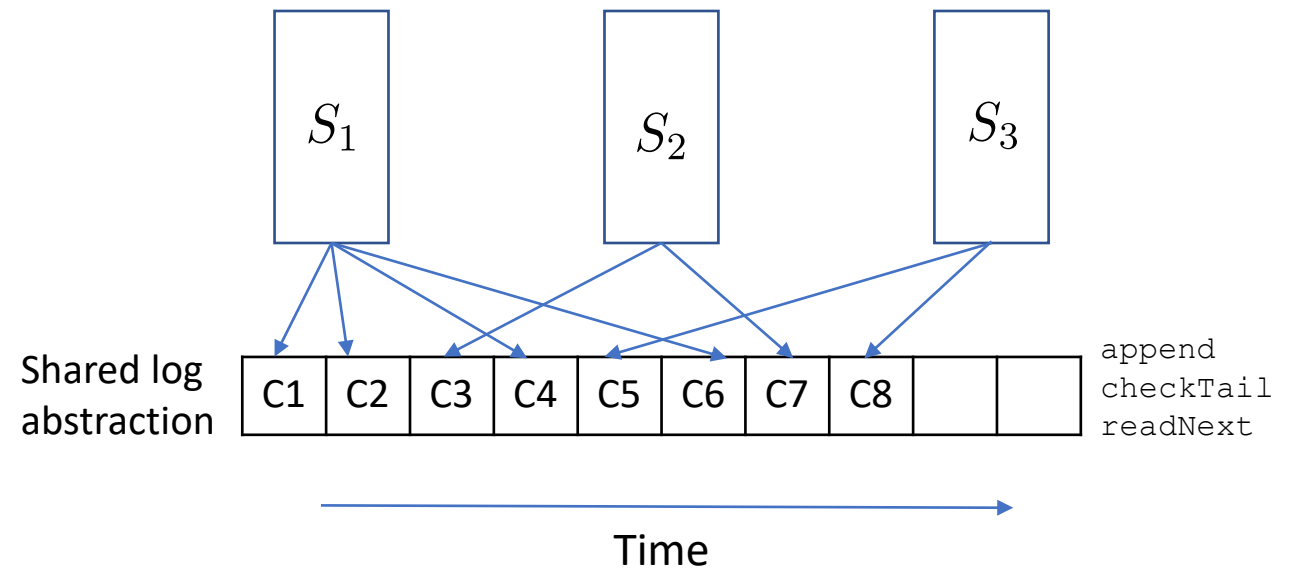
Shared log abstraction:

- append
- checkTail
- readNext

*achieved by*

Fault-tolerant consensus:

- Ordering protocol
- Leader election
- Membership changes

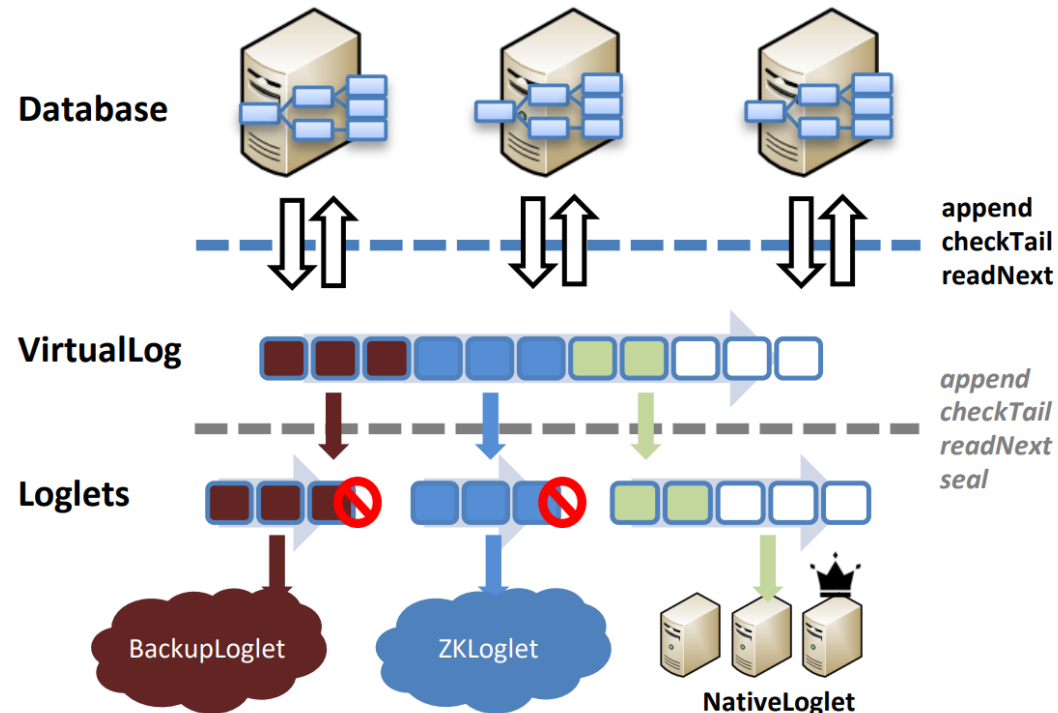


# Problem

- Consensus-based database replication systems are **monolithic, complex, and difficult to evolve**.
- Shared log a good abstraction, but implementations are:
  - *Difficult to deploy and operate*: **no support for upgrading and migrating applications without downtime**
  - *Difficult to develop*: monolithic consensus protocols with **coupled control plane and data plane**
    - Fault-tolerant consensus { [Data plane] Ordering protocol  
[Control plane] Leader election  
[Control plane] Membership changes

# Virtual consensus

- Layered approach:
  - **VirtualLog** for *control plane* (reconfiguration, leader election, membership changes)
  - **Loglets** for *data plane* (ordering commands and storing them durably)



# Key benefits

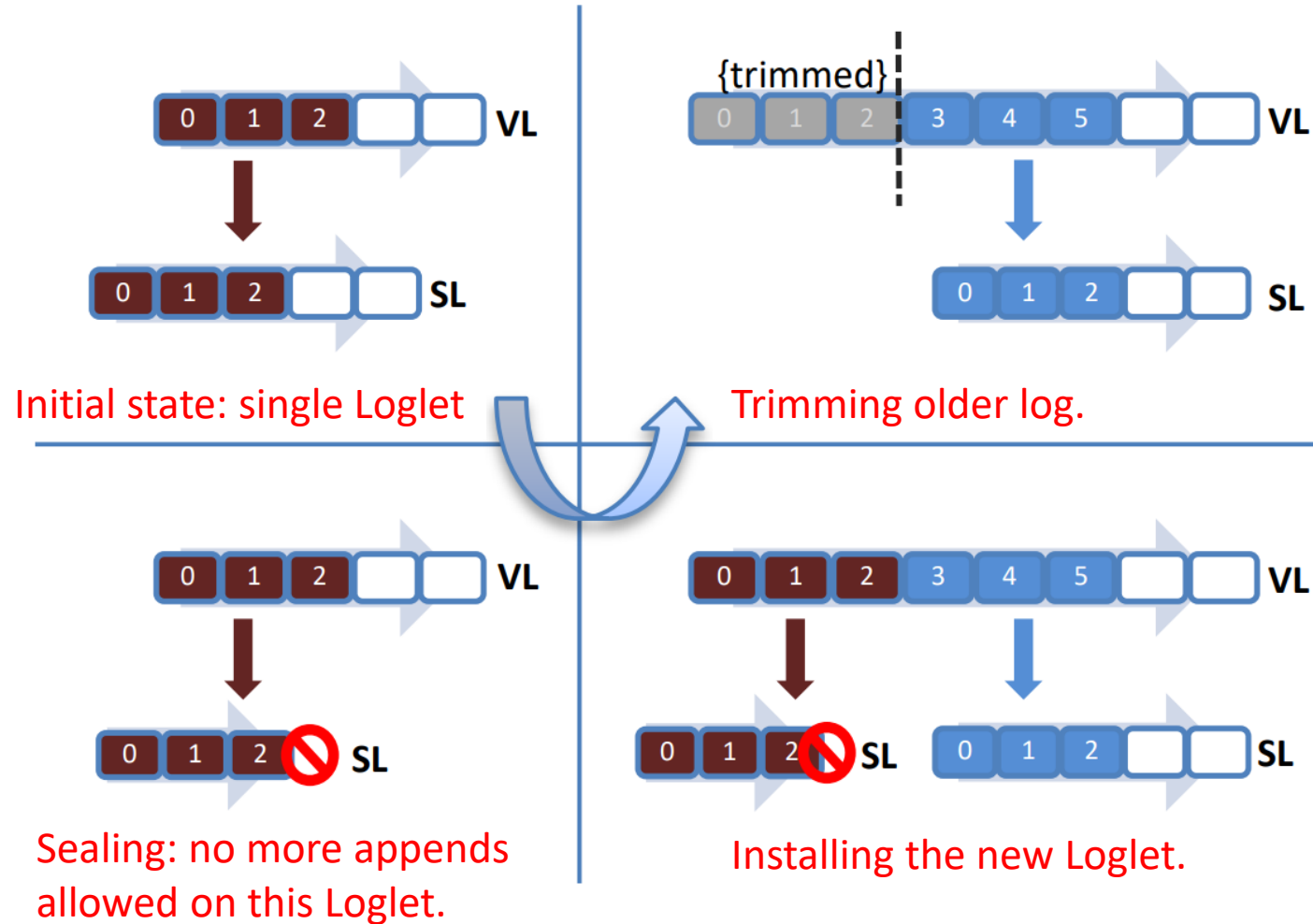
- Applications agnostic of VirtualLog's virtual nature, and see the traditional shared log API.
- Loglets do not need to implement a full **fault-tolerant** consensus protocol.
  - Can only implement a simple ordering protocol.\*
- VirtualLog's fault tolerant consensus does not need to be efficient, as it is triggered only during reconfiguration.
- Allows hot-swapping of Loglets as scalability requirements change.

\*: Loglets only need to implement a fault-tolerant 'seal' operation, which is theoretically weaker and much simpler to implement

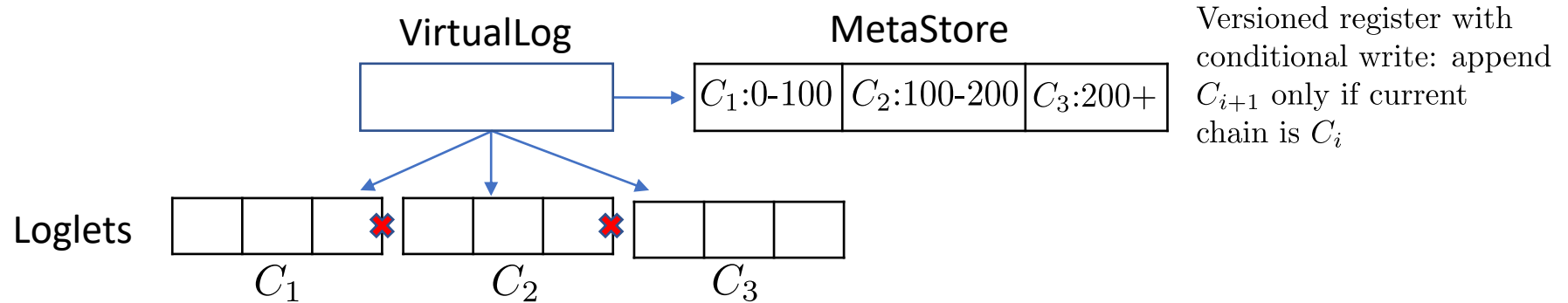
# Virtual consensus in Delos

- Delos: A database running on top of the VirtualLog API.
- *Simplified deployment and operation:*
  - Rapid initial development with rudimentary Loglet implementation
  - 10X improvement in latency by hot-swapping the Loglet implementation in production to a more efficient one.
  - Migrating older segments to a Loglet layered on cold storage.
- *Simplified development:*
  - *Loglets are simple to design:* Delos with a primary sequencer-based Loglet protocol.
  - *Loglets can be composed:* A StripedLoglet to achieve the following:
    - Double the failure threshold by rotating the sequencer.
    - Support >1M appends/sec by sharding.

# VirtualLog reconfiguration



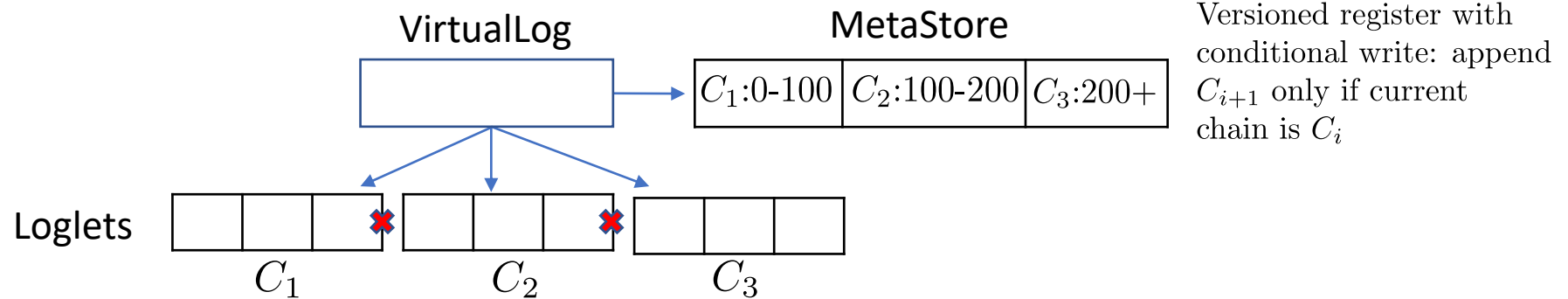
# VirtualLog reconfiguration protocol



1. Seal the old chain (seals are idempotent; appends disallowed post-sealing)
2. Install new chain on MetaStore (at most one winner in case of races)
3. Fetch new chain from MetaStore (in case someone else won the race)

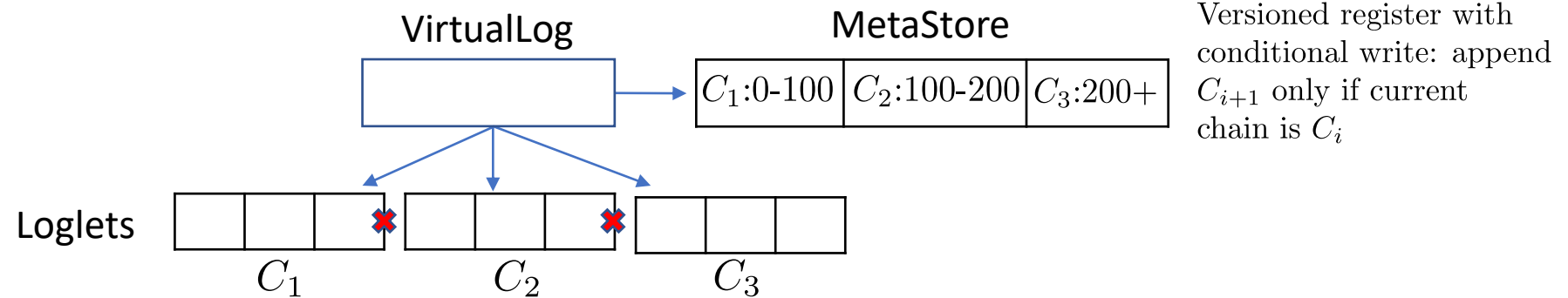


# Guarantees



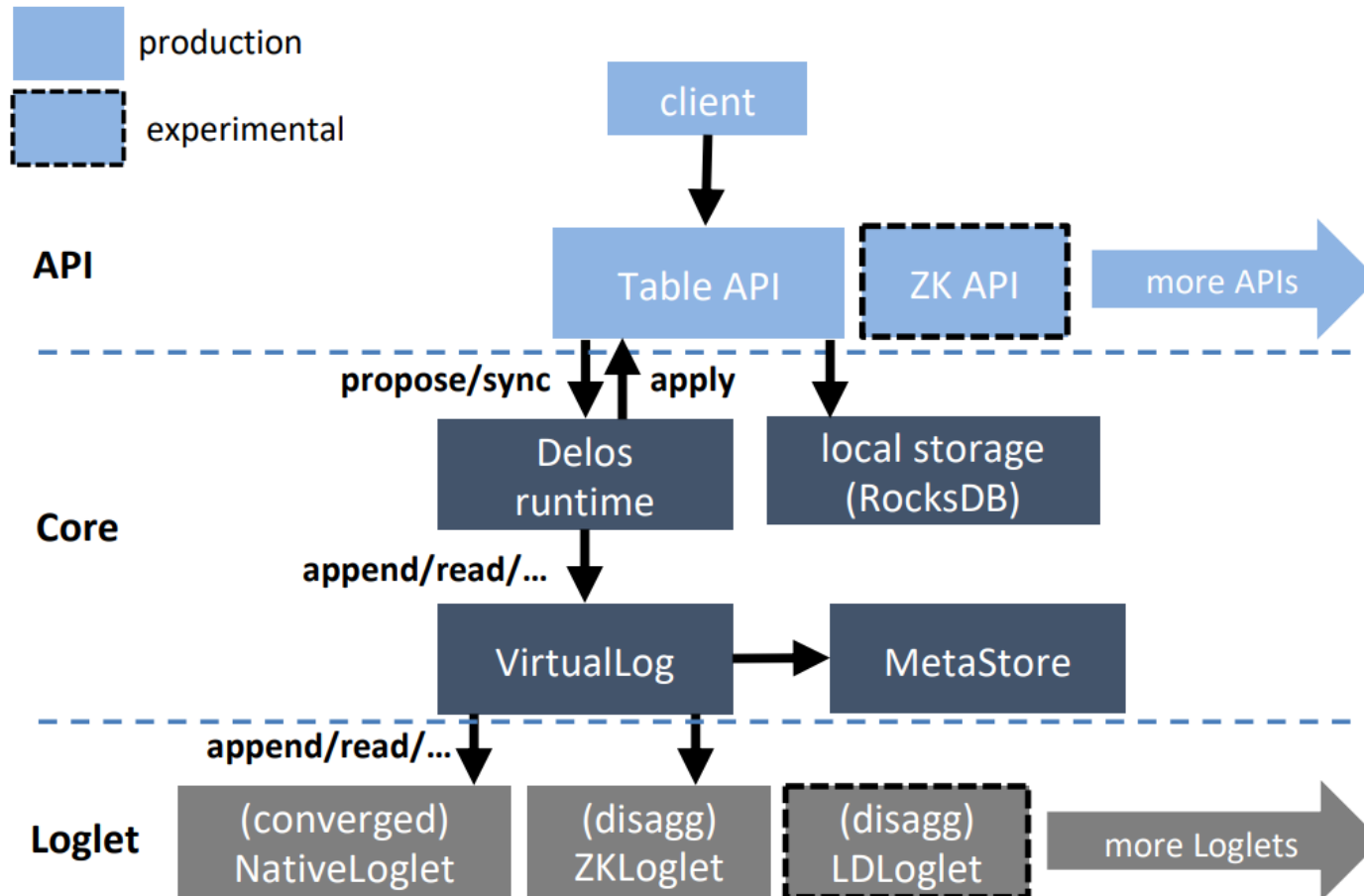
- **Concurrency:**
  - Multiple clients can seal concurrently. Seals are idempotent.
  - `checkTail` to a Loglet should return a *sealed* bit, indicating whether it is sealed or not. If working on a sealed chain, fetch the latest chain from MetaStore and try again.
- **Failure atomicity:**
  - If no new chain found after sealing, retry after some time and install a new chain with the same configuration as earlier.

# Requirements for components



1. MetaStore needs to implement the 'versioned register with conditional write', with a fault-tolerant consensus protocol like Paxos.
2. Loglets need to implement the sealing operation in a fault-tolerant way, via a *fault-tolerant atomic register*, which is weaker than consensus.
3. Loglets need to reject any append operations post-sealing.
4. Loglets need to return the *sealed* bit on the checkTail call.
5. Loglets need to detect failures and initiate the VirtualLog reconfiguration process.

# Delos – a database server over virtual consensus

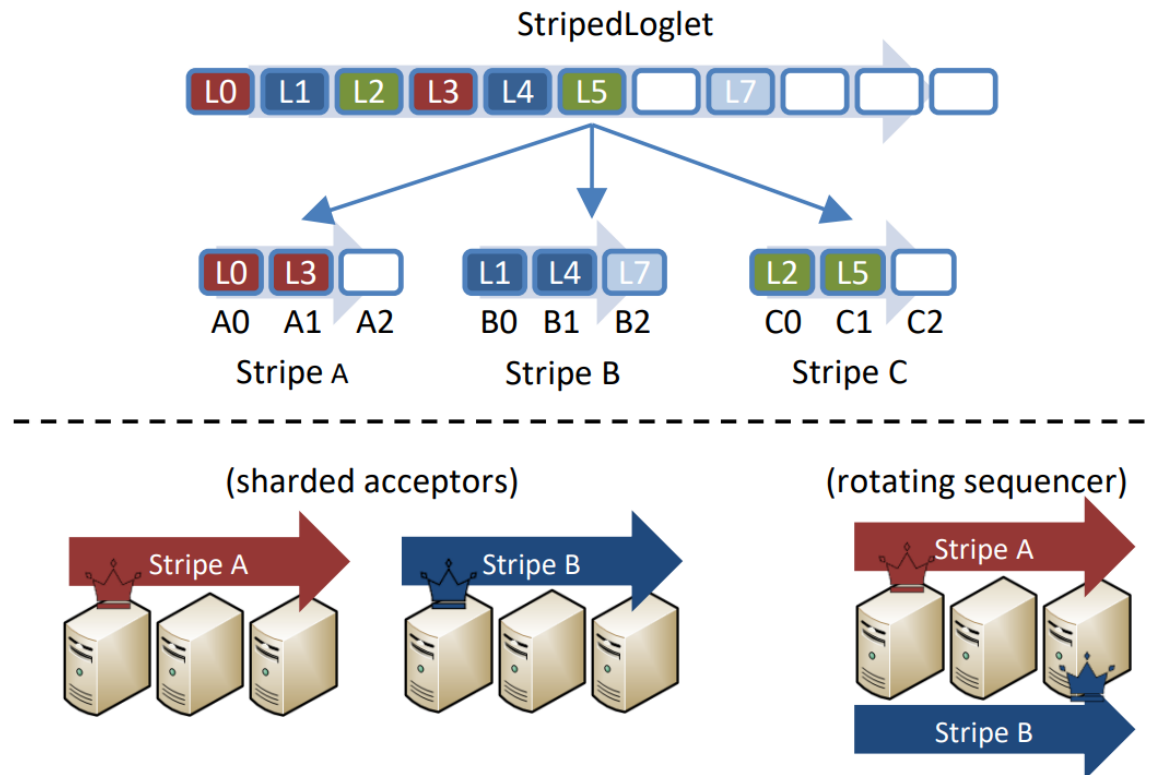


# Delos with NativeLogLets

- A primary sequencer-based Loglet protocol:
  - append calls contact a primary sequencer node and fail if the sequencer is down.
  - sealing requests each LogServer to set the seal bit.
  - checkTail implements a protocol to return a sealed status based on the sealed status of each LogServer.

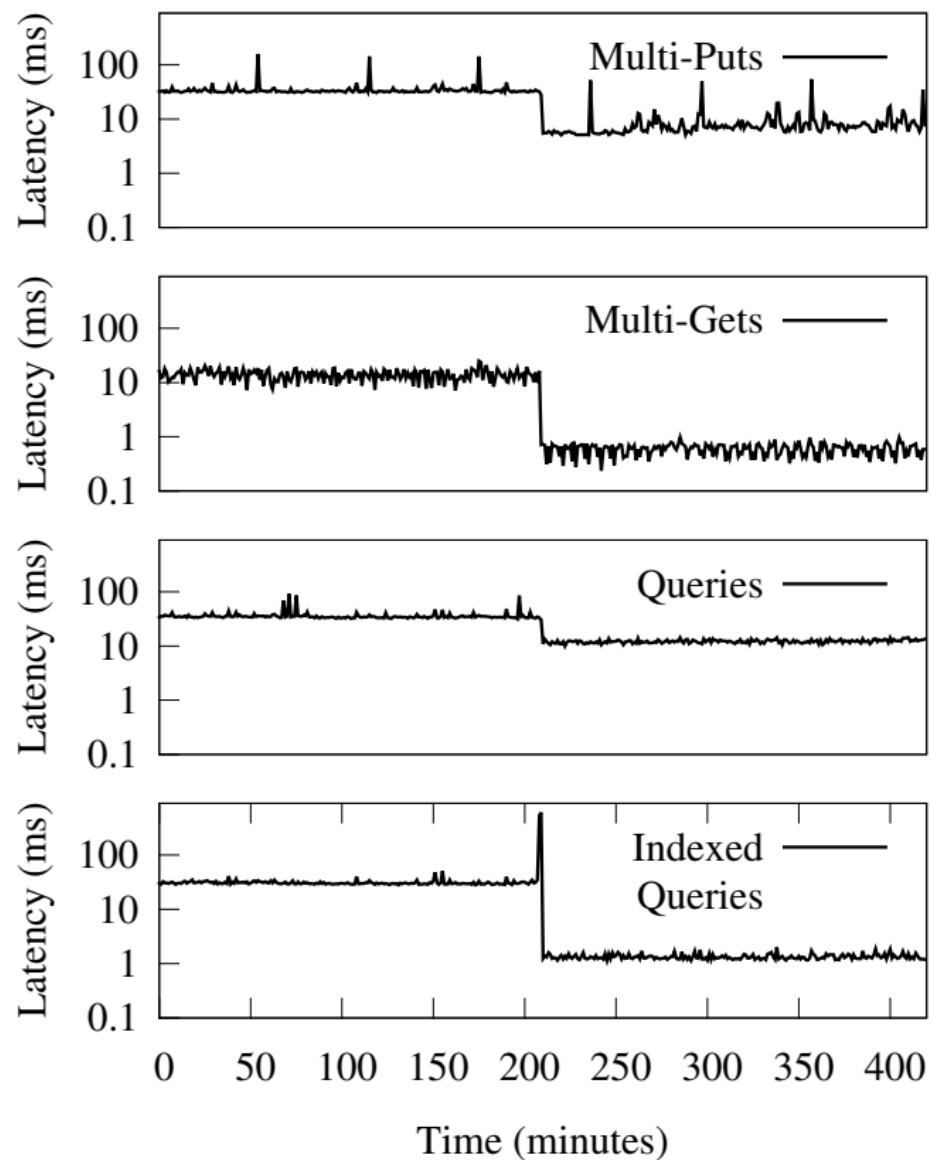
# Delos with StripedLogLets

- StripedLoglets compose simple Loglets to get special performance and robustness properties.

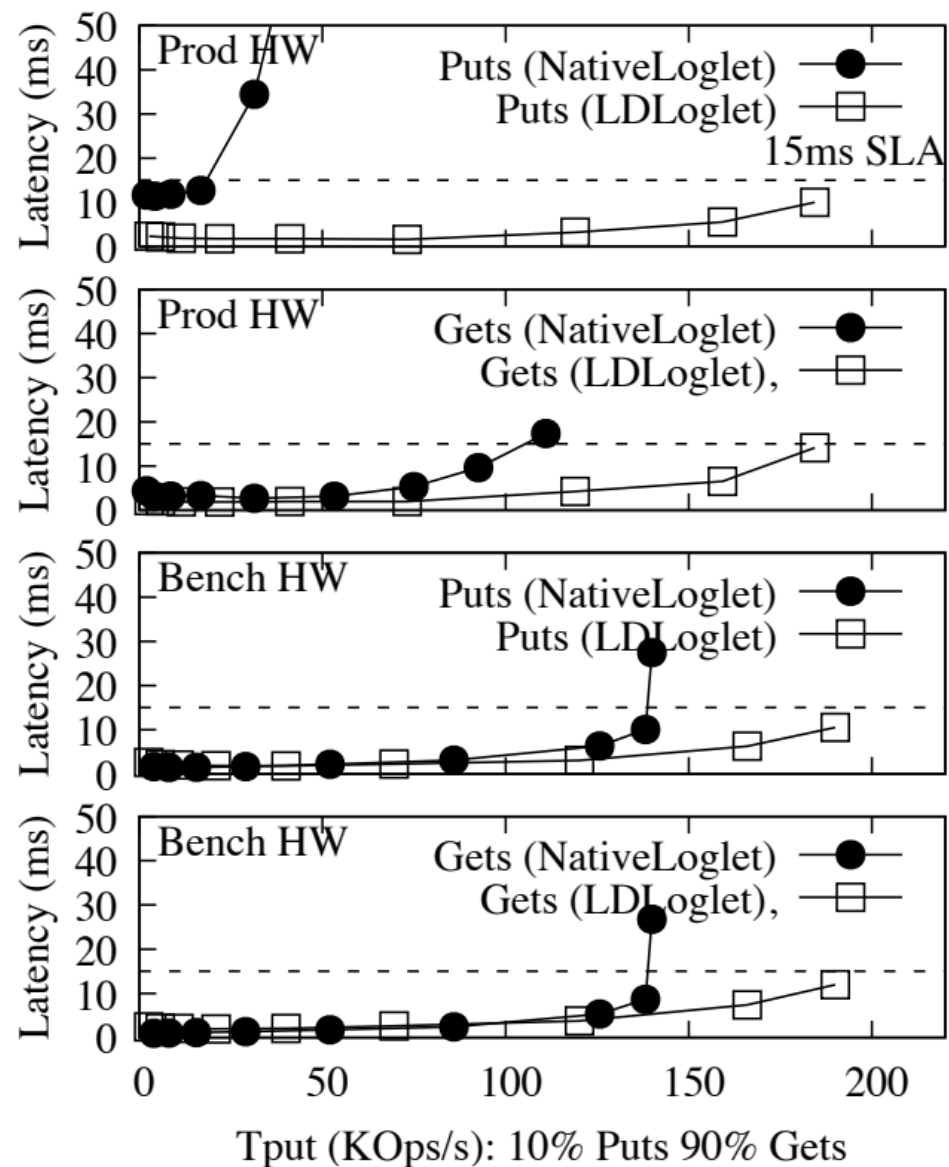


Benchmarks

# Latency during production switchover of the Loglet protocol

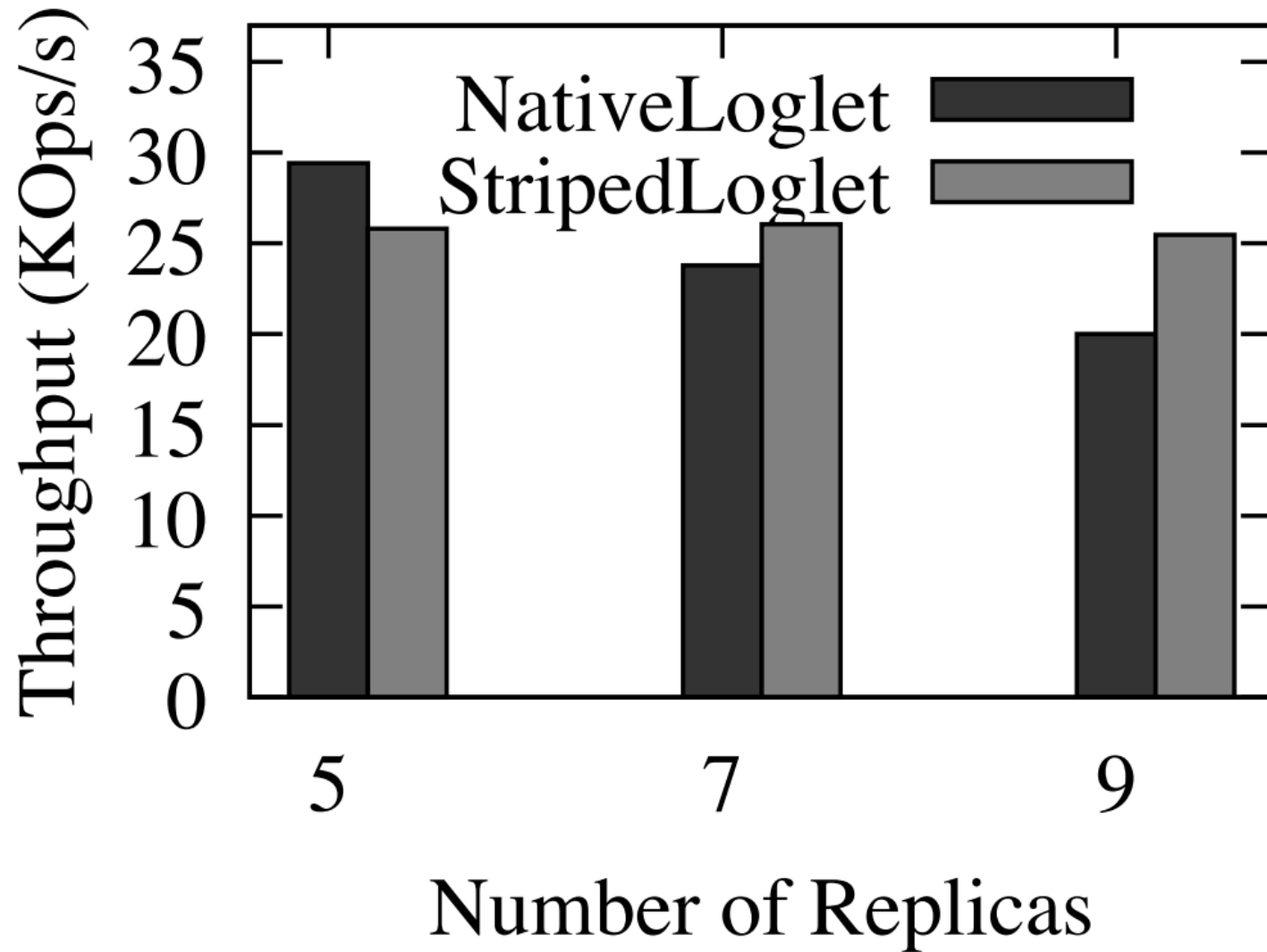


Throughput  
improvement  
via  
disaggregation

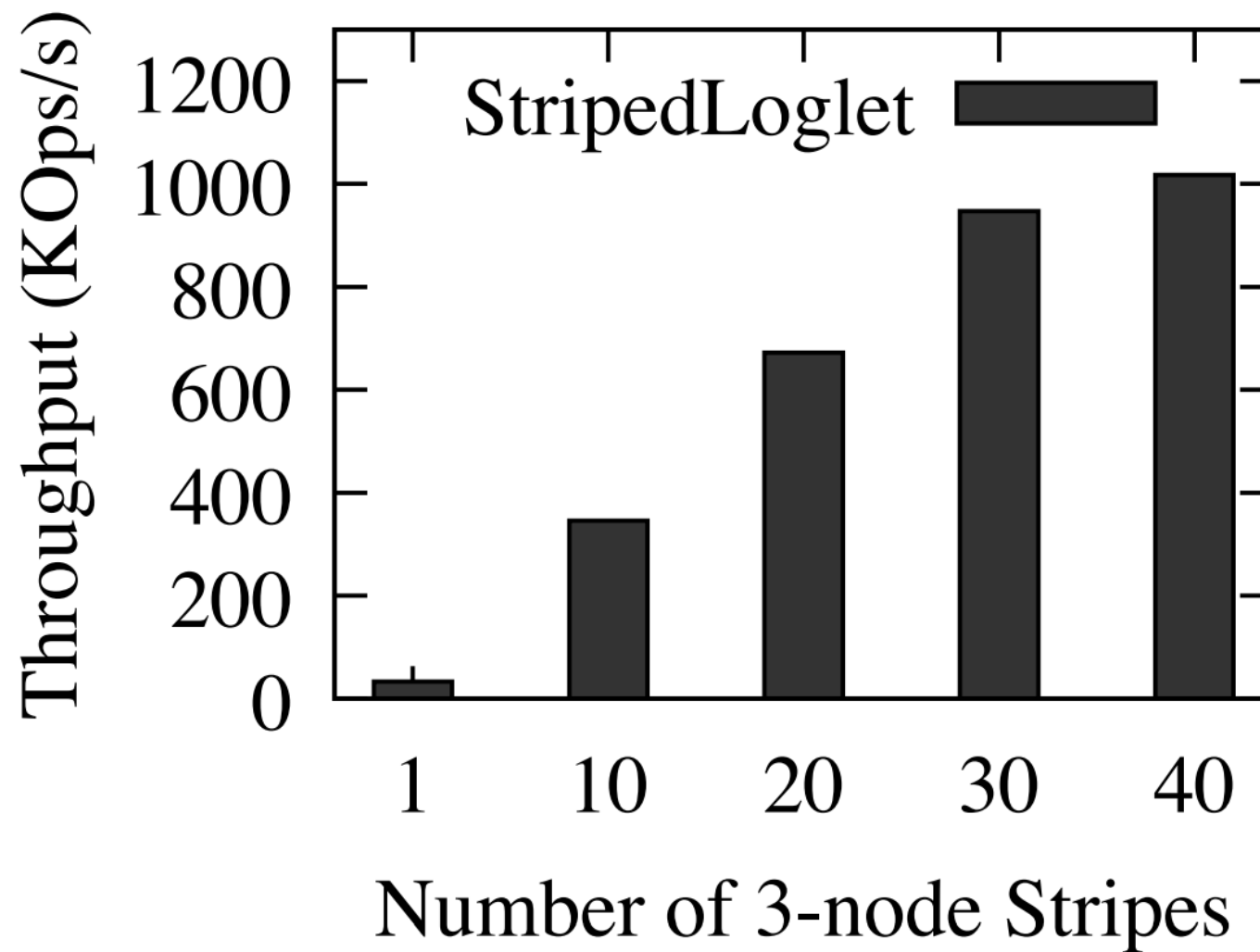




Removing  
sequencer  
bottlenecks

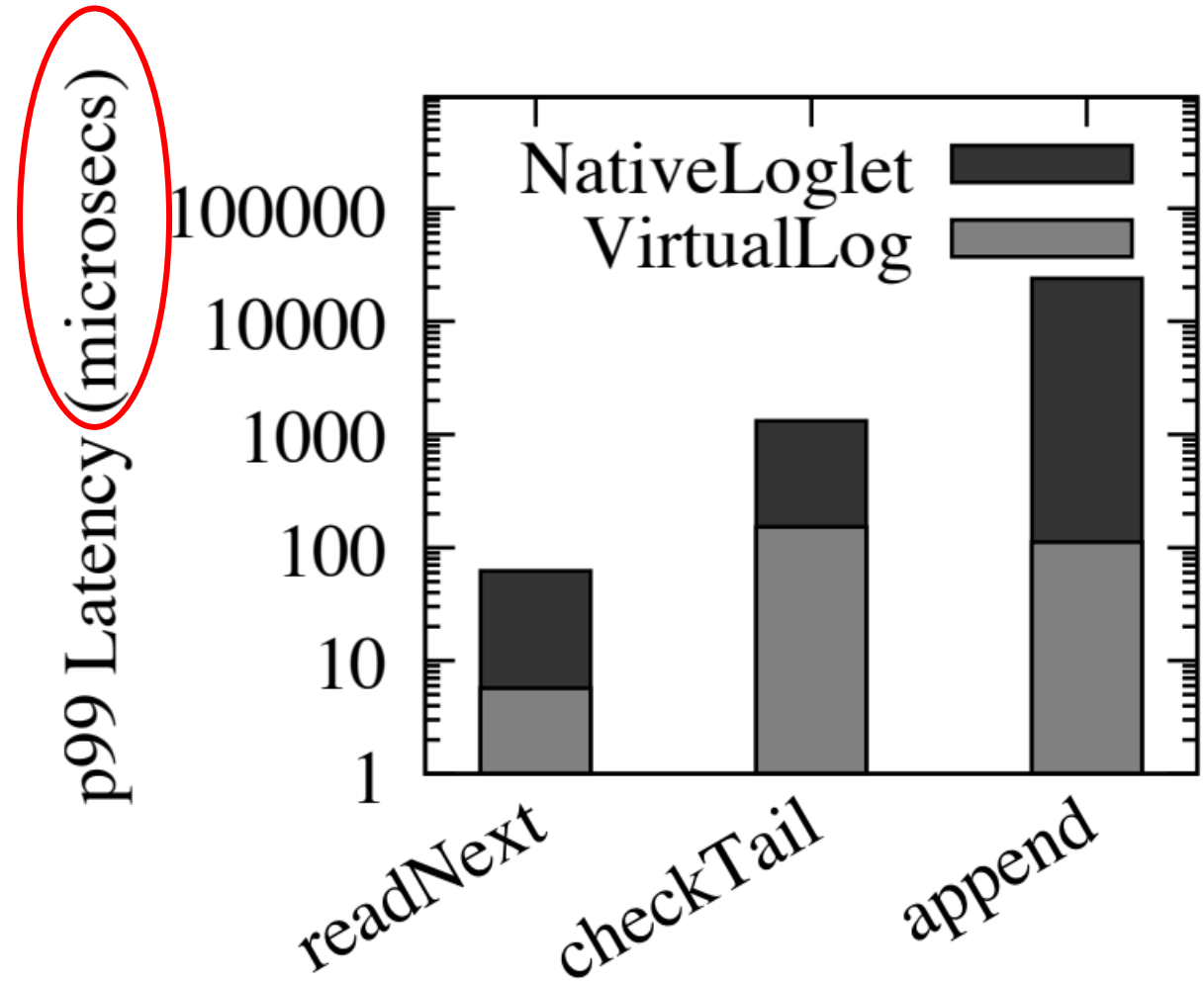


Horizontal  
scaling



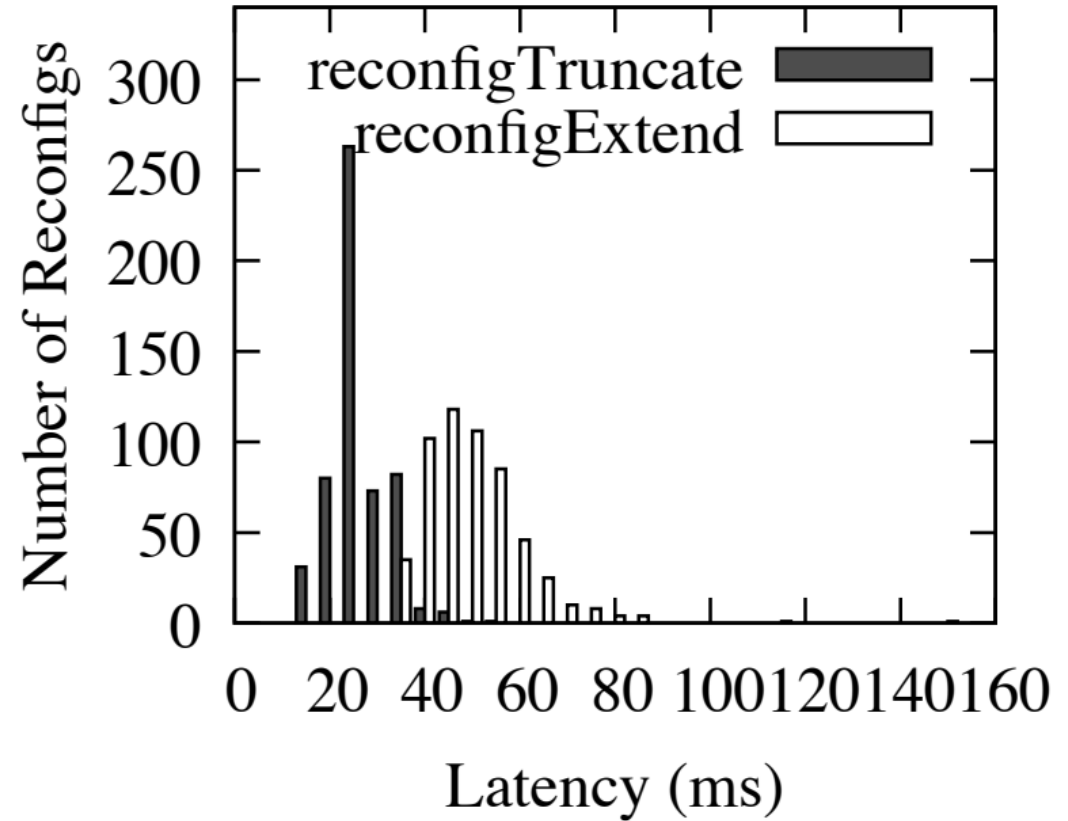
# The cost of virtualisation

Virtualisation  
latency



# The cost of virtualisation

Reconfiguration  
latency



# Comparison with ZooKeeper

